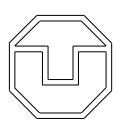
# Scalability in Computer Systems DNS/BIND as an example



## **Outline and Goal of Lecture**

#### Outline:

- Scalability: terminology, problems and basic approaches
- Names in Distributed Systems: purposes of naming, terminology
- Application of scalability approaches on name resolution

#### Goal:

understand some of the important principles how to build scalable systems

using DNS as an example

# **Definitions Scalability**

## Scalability

the ease with which a system or component can be modified to fit the problem area.

http://www.sei.cmu.edu/str/indexes/glossary/

Scalability [in telecommunication and software engineering] indicates the capability of a system to increase performance under an increased load when resources (typically hardware) are added

Wikipedia

# **Definitions Scalability**

#### Definition(Wang, Xu 98):

A computer system (HW + SW) is called scalable if it can scale up (improve its resources) to accommodate ever increasing performance and functionality demand and / or scale down (decrease resources) to reduce cost.

#### Dimensions of Scalability:

- Size (more CPUs)
- Other Resources (Memory)
- Software (Versions, better libs, etc.)
- Heterogeneity (different hardware / SW = portability)

# Scalability in Computer Systems

A system is described as scalable

if it remains effective when there is a significant increase in the number of resources and the number of users.

Coulouris, Dollimore, Kindberg: Distributed Systems

A system is scalable

if it works well for very large and very small numbers

# Another aspect of scalability:

Prepare for change in functionality

- software engineering
- choose sufficiently large logical resources
- provide hooks for extension

# Problems for Scalability in Distributed/Parallel Systems

Performance bottlenecks in ...

Failures of ...

Abuse of ...

- computers
- communication

# Principles to achieve Scalability ("RPC")

## **Partitioning**

split systems into parts that can operate independently to a large extent

#### Replication

provide several copies of components

- that are kept consistent eventually
- that can be used in case of failure of copies

# Locality (Caching)

maintain a copy of information

that is nearer, cheaper/faster to access than the original

# Some Challenges

#### Balance load:

keep load under reasonable threshold

- at each component
- in the communication subsystems

load balancing can be static or dynamic. Will study a detailed example for dynamic load balancing later(Mosix).

Minimize the delay induced by "RPC".

Prepare for change.

Information Dissemination.

Choose right degree of consistency.

## Names, Identifiers, Addresses

#### **Names**

- symbolic
- have a meaning for people

#### Identifiers

- identifies a component (uniquely)
- are used by programs

#### Addresses

- locates a component
- can change

## Name resolution

#### Name Resolution:

- map symbolic names to objects
- better: to a set of attributes such as identifiers, addresses, other names, security properties

#### Interfaces:

- Register (Name, attributes, ...)
- Lookup (Name) -> attributes

## Related

## Compilers

statically map names to addresses

## Dynamic libraries

dynamically remap addresses

## Port Mapper

map service to port

Name resolution is a form of dynamic mapping of pathnames to attributes.

## Observation

Many services, tools, ... provide their own name resolution

- file systems path names to I-Nodes
- login
- RPC systems (portmapper)

# Purpose of Directory Services

- integration of name services
- generic name service
- world-wide use of names

#### Today mostly used:

- email/web
- computer attributes (IP addresses)
- people attributes (certificates, ...)

## A Bit of History

## UUCP/MMDF (cum grano salis):

- ira!gmdzi!oldenburg!heinrich!user (path to destination)
- user@ira!heinrich%gmdzi
   (mixing identifiers and path information)

#### **ARPA-Net:**

- a single file: hosts.txt
- maintained at Network Information Center of SRI (Stanford)
- accessed via ftp
- TCP/IP in BSD Unix => chaos name collisions, consistency, load

DNS: Paul Mockapetries (84) ...

# More Terminology

## Name Space

set of names recognized by a name service

#### Context

unit for which a name can be mapped directly

#### Aliases

several names for one object

# **More Terminology**

## **Naming Domain**

Subtree in the hierarchy of DNS contexts

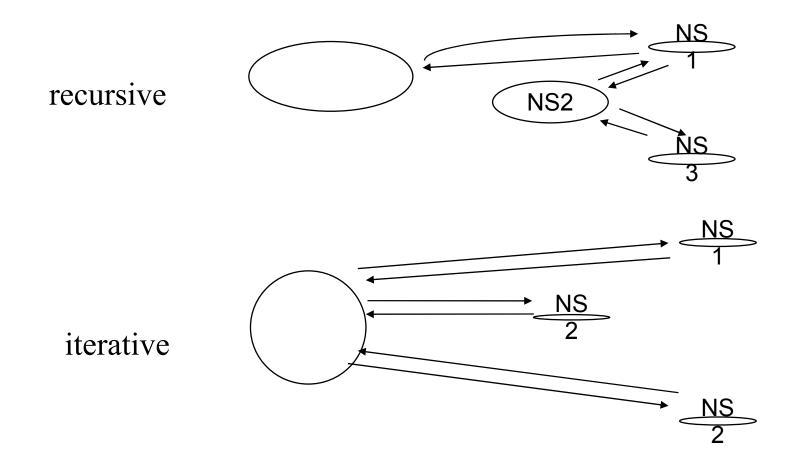
#### Zone

 (aka Zone of authority) Subset of a domain over which an authority has complete control. Subzones (starting at apices of a zone) can be delegated to other authorities.

#### Navigation

querying in a set of cooperating name spaces

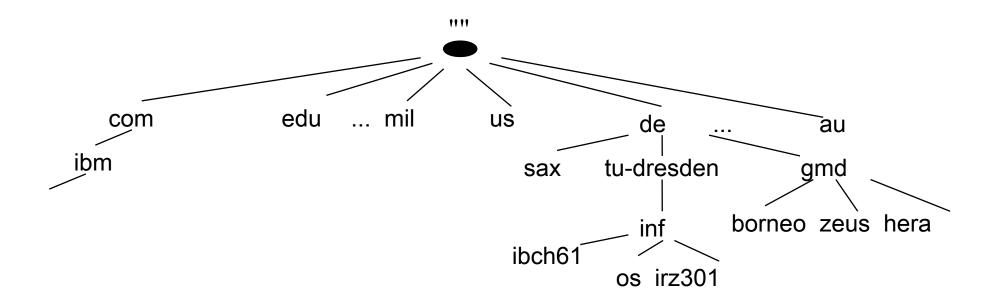
# **Basic Implementation Variants**



## Requirements / Properties

- arbitrarily large numbers
- arbitrary units of administration
- long living names, the higher in the hierarchy the longer
- high robustness
- restructuring of name spaces
- consistency
- efficiency

# **DNS Name Space**



## Examples

inf.tu-dresden.de Domain

os.inf.tu-dresden.de Computer

heidelberg.ibm.com Domain

ftp ftp.inf.tu-dresden.de

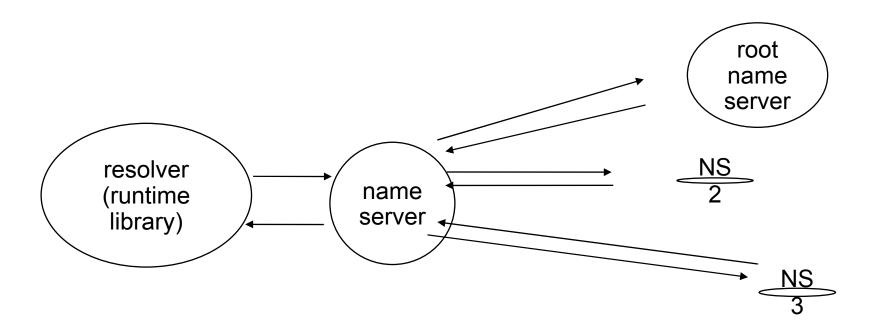
DNS: --> IP address: 141.76.2.3

ftp daemon: IP address, port 21

#### Properties:

- location independent
- not very deep

# Implementation Structure (BIND)



## Partitions: Zones

#### Zones:

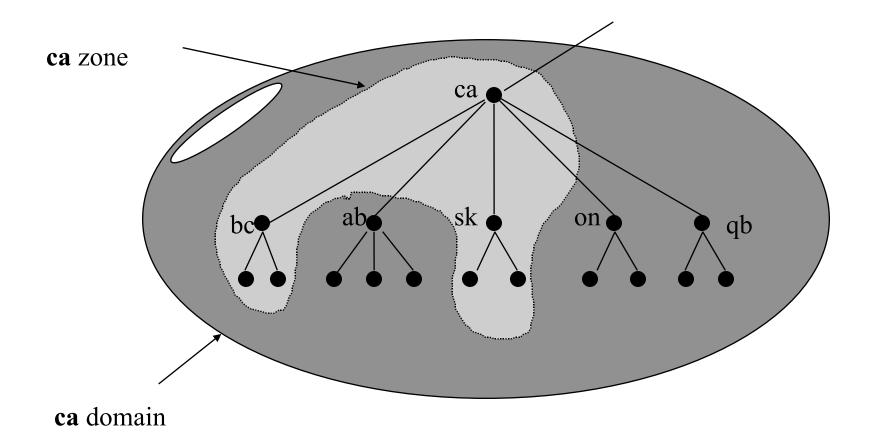
- administrative unit
- resolves all names within a zone recursively
- maps to names and addresses of name servers responsible for sub zones
- maintains management data

#### Name server:

process doing the name resolution for one zone

#### Resource records (RR):

key interface



# Replication

Currently 13 root name servers

each zone has at least

- one primary
- one secondary

name server

# Caching

each name server caches resource records time to live attribute

authoritative versus non-authoritative answers

## Resource Records

record type interpretation content

A address IPv4 address

AAAA address IPv6 address

NS name server DNS name

CNAME symbolic link DNS name of canonical name

SOA start of authority zone-specific properties

PTR IP reverse pointer DNS name

HINFO host info text description of host OS

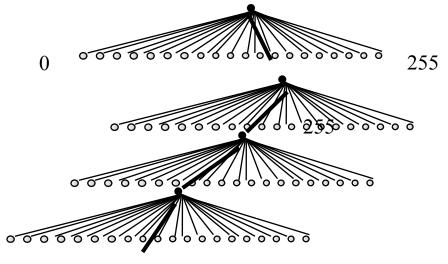
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## Reverse Resolution

Example

IP-Address: 141.76.48.97

=> DNS-Name: 97.48.76.141.in-addr.arpa



Xaver.os.inf.tu-dresden.de

# Summary: Scalability and DNS

## Good points:

- replication and caching work well
- over time, DNS scaled from small numbers to millions

#### **Bad Points:**

IP addresses too small

## Literature

Paul Albitz & Cricket Liu DNS and BIND O'Reilly & Associates, Inc.